

GSOL: A Confluence Checker for Haskell Rewrite Rules

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- GSOL: a tool to check the confluence of GHC rewrite rules
- The tool uses the Second-Order Laboratory (SOL) [1, 2] for checking the confluence
- GHC has the feature of user-defined **rewrite rules** [3]
- SOL is powerful in verifying the confluence of higher-order rewrite systems

Motivation

```
-- F.hs
module F where

{-# RULES
  "f/0" forall x. f x = 0 ;
  "f/1" forall x. f x = 1 ;
#-}

e = f 99

f :: Integer -> Integer
```

- "f/0" and "f/1" both match `f 99`.
- If GHC applies "f/0", `f 99` rewrites to 0.
- If GHC applies "f/1", `f 99` rewrites to 1.
- The rules "f/0" and "f/1" are *non-confluent*.
- GHC does not notice this non-confluence.

GSOL uses SOL to automate confluence checking of user-defined Haskell rewrite rules.

The Notion of Confluence



Confluence guarantees uniqueness of normal forms

- **Confluence (CR)** is when any two divergent computation paths are joinable.
- **Local Confluence (LCR)** is when all one step divergences are joinable.
- **Termination (SN)** is when there are no infinite rewrite steps.

GSOL uses **Newman's Lemma** to conclude CR by checking LCR and SN.

Implementation

- GSOL is implemented as a GHC Core plugin.



- GSOL's tasks
 1. Collect the Core rules
 2. Convert the Core rules to SOL rules
 3. Call SOL for checking:
 - SOL performs its checking functions,
 - SOL prints the results to standard output.

Second-Order Laboratory (SOL)

- SOL implements a formal framework of *second-order computation systems*.
- Second-order computation systems use the language of meta-terms:

$$t ::= x \quad | \quad x.t \quad | \quad f(t_1, \dots, t_n) \quad | \quad M[t_1, \dots, t_n].$$

- These forms are respectively variables, abstractions, and function terms, and meta-applications.
- The meta-terms have *second-order types*.
- Computation rules are pairs of meta-terms.

Stand-alone command

```
gsol [PROPERTY] FILENAME
```

- PROPERTY
 - `cri` local confluence checking
 - `sn` termination checking
- FILENAME path to the Haskell program file

Web Interface

<http://solweb.mydns.jp/>

Example: Arrow

- Control.Arrow is a library in GHC
- It offers a way to programming with various computational effects
- It is implemented with the type class:

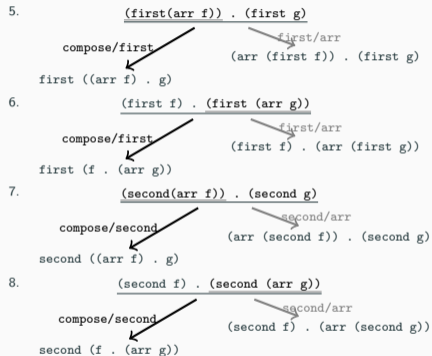
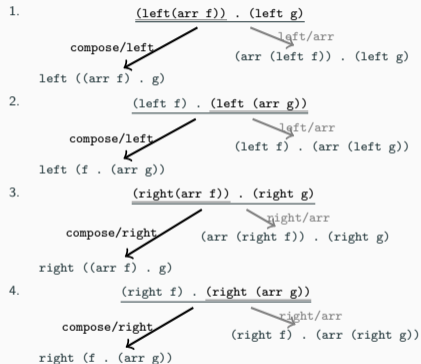
```
class Category a => Arrow a where
  arr      :: (b -> c) -> a b c
  first   :: a b c -> a (b,d) (c,d)
  ...
```


Rewrite rules in Control.Arrow

```
{-# RULES
"compose/arr"    forall f g. (arr f) . (arr g)    = arr (f . g)
"first/arr"      forall f . first (arr f)        = arr (first f)
"second/arr"     forall f . second (arr f)       = arr (second f)
"compose/first" forall f g. (first f) . (first g) = first (f . g)
"product/arr"    forall f g. arr f *** arr g     = arr (f *** g)
...
#-}
```

- These rewrite rules are described in Control.Arrow
- NB: The laws assume $a = (->)$

8 critical pairs in the Arrow Laws



A closer look at a non-joinable critical pair

```
"compose/first" forall f g. (first f) . (first g) = first (f . g)
"first/arr"      forall f'. first (arr f') = arr (first f')
```

- Since `(first f)` is unifiable with `first (arr f')` by $f \mapsto \text{arr } f'$ it generates a critical pair:

$$\frac{\text{first}(\text{arr}(f')) . (\text{first } g)}{\text{compose/first} \quad \text{first } (\text{arr}(f')) . g \neq \text{arr } (\text{first}(f')) . \text{first}(g) \quad \text{first/arr}}$$

- Since these are normal forms, it shows **non-confluence** of the rules in `Control.Arrow`.
- Useful information to redesign the rules.

Future work

- Better handling of types in the translation from Core rules to SOL rules
 - GSOL ignore type constraints
 - Enough for checking critical pairs, but termination checking is weaker

Summary

- GSOL is a tool to check the confluence of Haskell rewrite rules
- GSOL uses the Second-Order Laboratory (SOL) for confluence checking
- Web interface <http://solweb.mydns.jp/>
- Draft paper <http://solweb.mydns.jp/gsol-draft.pdf> [4].

- [1] M. Hamana, “How to prove decidability of equational theories with second-order computation analyser SOL,” *Journal of Functional Programming*, vol. 29, no. e20, 2019.
- [2] M. Hamana, T. Abe, and K. Kikuchi, “Polymorphic computation systems: Theory and practice of confluence with call-by-value,” *Science of Computer Programming*, vol. 187, no. 102322, 2020.
- [3] S. P. Jones, A. Tolmach, and T. Hoare, “Playing by the rules: rewriting as a practical optimisation technique in GHC,” in *Haskell Workshop 2001*, 2001.
- [4] Y. F. Date and M. Hamana, “GSOL: A confluence checker for Haskell rewrite rules.” <http://solweb.mydns.jp/gsol-draft.pdf>.